

**Exercise 1** LRU vs FIFO

*3+3+1 Points*

Recall the two online paging algorithms, *least recently used* (LRU), and *first-in-first-out* (FIFO). The first algorithm evicts, in case of a cache miss, the page whose last request time is the earliest. The second algorithm evicts the page whose insertion time is the earliest. Suppose the cache size is  $K = 3$ . You may assume in all cases that the cache is initially empty.

- (a) Give an example sequence  $R_1$  where the ratio  $LRU(R_1)/FIFO(R_1)$  is arbitrarily close to 3.
- (b) Give an example sequence  $R_2$  where the ratio  $FIFO(R_2)/LRU(R_2)$  is arbitrarily close to 2.
- (c) Can these ratios be larger than 3?

**Exercise 2** Least frequently used

*3 Points*

Show that the online paging algorithm *least frequently used* (LFU) can not be  $f(K)$ -competitive for any function  $f$ , where  $K$  is the cache size. Recall that LFU evicts, in case of a cache miss, the page that has been requested the fewest number of times since insertion, among those pages that are currently in the cache.

**Exercise 3** Changing the cache size

*4+0 Points*

Let  $LRU(R)$  denote the cost of the LRU algorithm for request sequence  $R$  with cache size  $K$ , and let  $LRU'(R)$  denote the cost of LRU for request sequence  $R$  with cache size  $K + 1$ . In both cases we start with an initially empty cache.

- (a) Show that  $LRU'(R) \leq LRU(R)$ , or in other words, increasing the cache size cannot increase the cost.
- (b) Perhaps surprisingly, this is not true for every algorithm, e.g. for FIFO one can construct an example sequence where increasing the cache size *increases* the cost. Study the example in the Wikipedia article “Bélády’s anomaly”.
- (c) Bonus (+4p): Check whether this phenomenon can happen for the LFU, LIFO, and LFD algorithms.

**Exercise 4** Programming exercise

Please submit the programming exercise report and source code by June 29th.

*Total: 14 points. Have fun with the solutions!*