

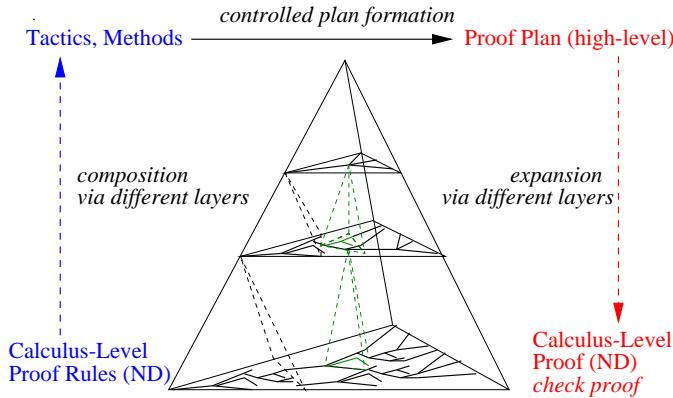


Proof Transformation and Expansion with a Parameterizable Inference Machine

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Hierarchical Proof Data Structure



- * Proof creation at abstract level with tactics and methods
- * Expansion in a three dimensional data structure

Procedural Expansion

Expansion of tactic PUSHNEG:

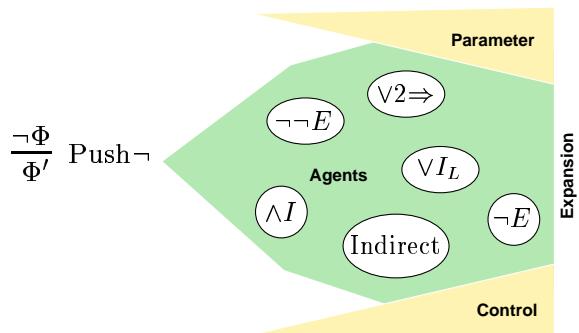
$\neg(A \vee B)$	Push \neg	(I.0), t ₀ $\vdash \neg(A \vee B)$	Hyp
		(I.1), t ₁ $\vdash B$	Hyp
		(I.2), t ₂ $\vdash A$	Hyp
		(I.3), t ₀ , t ₂ $\vdash [A \vee B]$	$\vee I_L, 1, 2$
		(I.4), t ₀ , t ₂ $\vdash \perp$	$\neg E, 1, 0, 1, 3$
		(I.5), t ₀ $\vdash \neg A$	$\neg I, 4$
		(I.6), t ₀ , t ₁ $\vdash [A \vee B]$	$\vee I_R, 1, 1$
		(I.7), t ₀ , t ₁ $\vdash \perp$	$\neg E, 1, 0, 1, 6$
		(I.8), t ₀ $\vdash \neg B$	$\neg I, 7$
		(I.9), t ₀ $\vdash (\neg A \wedge \neg B)$	$\wedge I, 1, 5, 1, 8$

$\neg(A \wedge B)$	Push \neg	(I.0), t ₀ $\vdash \neg(A \wedge B)$	Hyp
		(I.1), t ₁ $\vdash \neg(A \wedge \neg B)$	Hyp
		(I.2), t ₂ $\vdash \neg B$	Push $\neg, 1, 1$
		(I.3), t ₃ $\vdash \neg \neg B$	$\neg E, 1, 2$
		(I.4), t ₁ $\vdash B$	$\neg \neg E, 1, 3$
		(I.5), t ₁ $\vdash \neg \neg A$	$\wedge E, 1, 2$
		(I.6), t ₁ $\vdash A$	$\neg E, 1, 5$
		(I.7), t ₁ $\vdash [A \wedge B]$	$\wedge I, 1, 4, 6$
		(I.8), t ₁ $\vdash \perp$	$\neg E, 1, 0, 1, 7$
		(I.9), t ₀ $\vdash (\neg A \vee \neg B)$	Indirect L ₈

$\neg(A \Rightarrow B)$	Push \neg	(I.0), t ₀ $\vdash \neg(A \Rightarrow B)$	Hyp
		(I.1), t ₁ $\vdash \neg(A \wedge \neg B)$	Hyp
		(I.2), t ₂ $\vdash \neg B$	Hyp
		(I.3), t ₃ $\vdash \neg \neg B$	Push $\neg, 1, 1$
		(I.4), t ₁ $\vdash [A \vee \neg B]$	$\vee I_L, 1, 2$
		(I.5), t ₁ $\vdash (\neg A \vee \neg B)$	$\vee I_R, 1, 1$
		(I.6), t ₁ $\vdash B$	$\neg E, 1, 3$
		(I.7), t ₁ $\vdash \neg (A \vee B)$	$\neg E, 1, 6$
		(I.8), t ₀ , t ₁ , t ₂ $\vdash (A \Rightarrow B)$	$\vee 2, 1, 7$
		(I.9), t ₀ , t ₁ , t ₂ $\vdash (A \Rightarrow B)$	$\vee 2, 1, 5$
		(I.10), t ₀ , t ₁ , t ₂ $\vdash (A \Rightarrow B)$	$\vee E, 1, 4, 1, 8, 1, 9$
		(I.11), t ₀ , t ₁ $\vdash \perp$	$\neg E, 1, 0, 1, 10$
		(I.12), t ₀ $\vdash (A \wedge \neg B)$	Indirect L ₁₁

- * Justification: tactic/method + parameters (for expansion)
- * Expansion: hardwired programming code
- * Modifications of a tactic results in reimplementing
 - the expansion procedure of the tactic itself
 - other expansion procedures employing this tactic

Expansion with Ω Ants

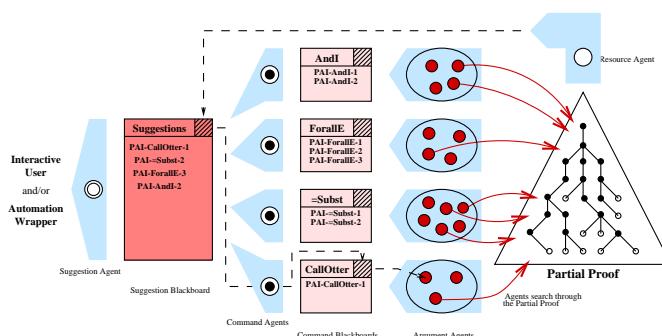


- * Employing the Ω Ants mechanism with
 - * Inferences
 - * Control Information
- creates expansion by proof search

Discussion:

- + Not hardwired, reduced maintenance
- + Flexible and adaptable
- ? Non-determinism in:
 - Proof search
 - Sub-proof structure (abstraction level)
- ? Limits with respect to procedural control information

Agent Based Deduction with Ω Ants



- * Ω Ants is a parameterizable inference machine
- * Agents and heuristics can be defined at run time
- * Incorporation of control information by pre-instantiation of blackboards and heuristics