# Formalising Category Theory by Automating Free Logic in Higher-Order Logic

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#### Introduction

Free logic extends classical logic by means to address non-existence of objects and to handle undefinedness in a suitable way. Reasoning about free logics can be realised, similar to other non-classical logics, by utilising semantical embeddings in classical higher-order logic [1]. Such an embedding has been employed for the axiomatisation of category theory, a theory, that depends on partiality and undefinedness and therefore significantly benefits from free logic in its formalisation. Parts of the category theory book *Categories*, *Allegories* by P. J. Freyd and A. Scedrov [4] have been formalised using an embedding of free logic in higher-order logic implemented in the mathematical proof assistant Isabelle/HOL [5]. The axiom system as presented in the book has been found to suffer from a constricted inconsistency [2], and hence an alternative axiom system provided by D. S. Scott [6] has been used in our formalisations.

object  $\star \in D \notin E$ . A graphical illustration of this notion of free logic is presented in Figure 2.

D: nicht-existente Objekte

The lemma shown in Figure 4, the equivalence of two functor definitions, could not be proved when formalized in a naive way. It is easy to show, that the first definition follows from the second one, but when assuming the first one and trying to prove the second one, the provers failed. Especially the third line of the second definition, the one with the directed equality, is not provable from the first definition. However, there exist some additional conditions, which can be added to make it derivable. However, it is unclear if Freyd and Scedrov did intend such conditions. Another issue to reconsider is the use of *if* and *iff* in the definitions. It has to be clarified if there is an intentional purpose for this distinction.



Figure 2: Graphical illustration of a domain and its subdomain

# **First Formalisations**

Category theory after Freyd and Scedrov [4] introduces the source operator  $(\Box X)$ , the target operator  $(Y\Box)$ , as well as partial composition  $X \cdot Y$  of two morphisms X and Y. Freyd and Scedrov assume Kleene equality (=) in most cases, and a directed Kleene equality ( $\geq$ ) in some special cases. Most of the formalised lemmas could be verified very easily by the theorem provers, for example, the lemma in Figure 3; only for some statements the provers failed.

## **Related Work**

Automating Free Logic in Isabelle/HOL [3] and Axiomatising Category Theory in Free Logic [2] by C. Benzmüller and D. Scott.

### **Results and Conclusion**

The first chapter of *Categories*, *Allegories* has been formalised. There are many chapters to go, but with the help of a powerful proof assistant and an appropriate embedding techniques a complete formalisation of the book may be

**Figure 1:** Front page of *Categories, Allegories* [4]

#### Free Logic

Free logic is a logic free of any existence presumptions. While preserving the original quantification over a specific domain, terms may now denote undefined/non-existing objects. Free logic after Scott [7] distinguishes between a raw domain D and a particular subdomain E of D. D contains possibly nonexisting objects while E holds only the existing entities. Free variables range over domain D and quantified variables only over domain E. Undefinedness is symbolized by a unique 1.13.  $\Box(\Box x) = \Box x$  because  $\Box(\Box x) = \Box((\Box x)\Box) = (\Box x)\Box = \Box x.$ Similarly  $(x\Box)\Box = x\Box$ .

Figure 3: Lemma 1.13. of [4]: Some equations

**1.18.** Given categories A and B, a function  $F: \mathbf{A} \rightarrow \mathbf{B}$  is a FUNCTOR if

> implies  $\Box(Fx) = Fy$ ,  $\Box x = y$ implies  $(Fx)\Box = Fy$ ,  $x\Box = y$ implies (Fx)(Fy) = Fz. xy = z

Equivalently, F is a functor iff

 $F(\Box x) = \Box(Fx) ,$ 

 $F(x\Box) = (Fx)\Box,$ 

 $F(xy) \simeq (Fx)(Fy)$ .

Figure 4: Lemma 1.18. of [4]: Equivalence of two functor definitions

in reach.

### References

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Leibniz: *Calculemus!* 

Computational Metaphysics is a interdisciplinary lecture course designed for advanced students of computer science, mathematics and philosophy. The main objective of the course is to teach the students how modern proof assistants based on expressive higher-order logic support the formal analysis of rational arguments in philosophy (and beyond). In our first course in Summer 2016 the focus has been on ontological arguments for the existence of God. However, some students picked formalisation projects also from other areas (including maths).

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